

Robust validation of network designs under uncertain demands and failures

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Validating network design

- Network design today is ad-hoc, and validating design is usually an afterthought
 - Contrast: Tools for chip and software industry a \$10B business [Mckeown, 2012]
- Much progress on verification of network data plane (e.g., reachability, security policy)
 - HSA, Veriflow, Batfish, NoD, etc.
- **Our goal:** Validating **quantitative** network properties
 - Formal approach to guarantee network performance (e.g., bandwidth, link utilization)
 - Under diverse failure/traffic scenarios
 - Use the formal approach to inform network design

Why is network validation hard? (1)

- **Scenarios of interest are too many**
 - Exponentially many failure scenarios [Wang et al., Sigcomm '10, Liu et al., Sigcomm '14]
 - E.g., All possible simultaneous f link failures
 - All possible traffic demands — non-enumerable

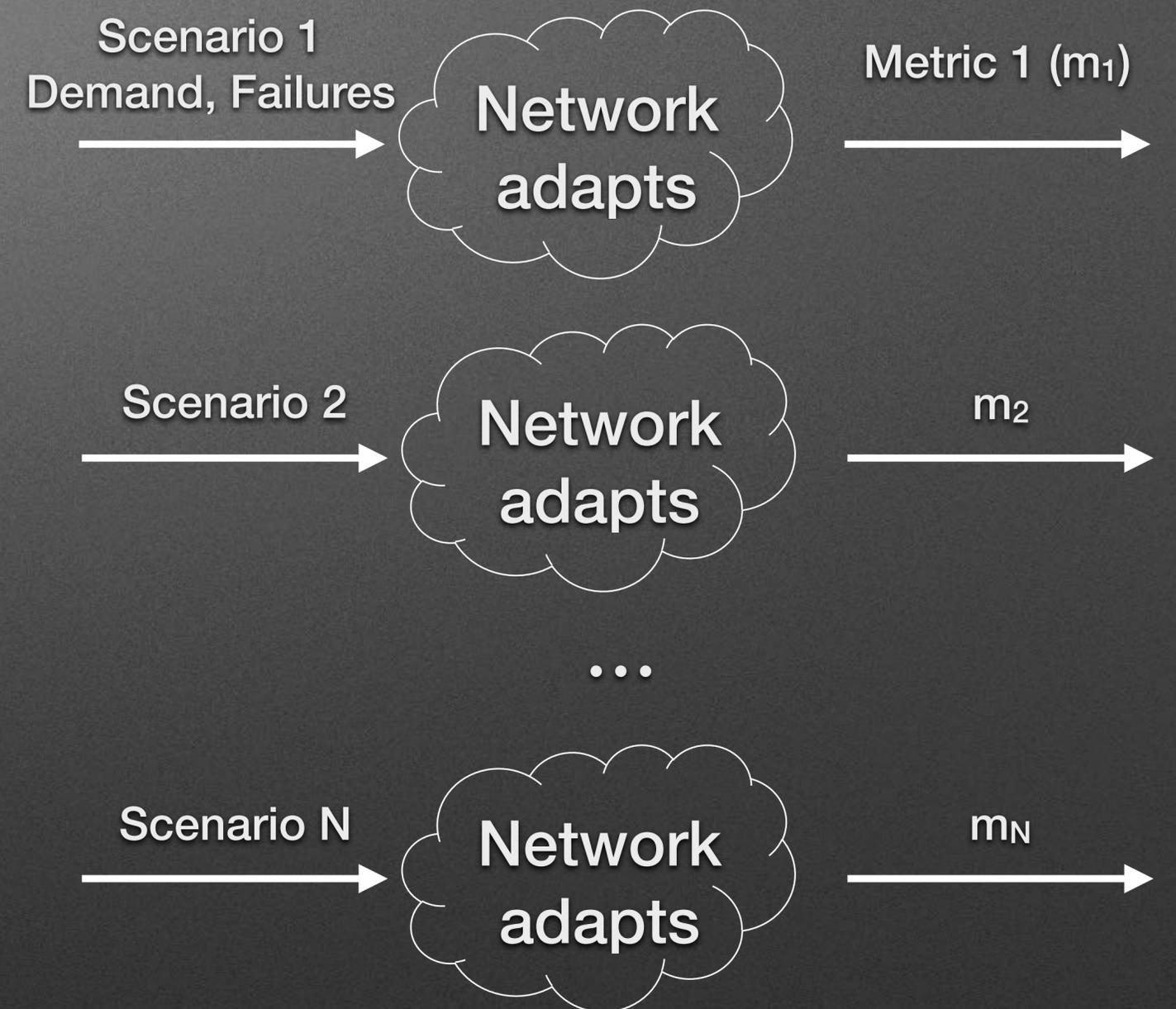
Why is network validation hard? (2)

- **Adaptation makes the problem intractable**
- Networks increasingly agile and flexible in adaptation
 - E.g., SDNs and NFVs
- Tools exist to bound worst case performance
 - E.g., robust optimization, and oblivious routing [Applegate et al., Sigcomm '03]
 - Assume networks do not adapt, or consider limited forms of adaptation to make problem tractable



Our work

- General framework for network validation
 - Find the **worst** performance of the network across **all** scenarios assuming network can adapt in **best** fashion for each scenario
- Handles intractable problems drawing on cutting-edge optimization technique
- Applies to **network synthesis**



Worst performance = $\max\{m_1, m_2, \dots, m_N\}$
Less is better

Example: Failure validation

Uncertainty Set

- All f or fewer link failures

Adaptations

- Flexible rerouting (multi-commodity flow)

Performance metric

- Utilization of most congested link

Problem:

- Given up to f links may simultaneously fail, what is the worst-case utilization of any link across all failure scenarios?

Wide applicability of framework

Uncertainty Set

- All f or fewer link failures
- Shared risk link group
- Weighted averages of historical demands

Adaptations

- Flexible rerouting (multi-commodity flow)
- Rerouting constrained to pre-selected tunnels
- Constrain with middlebox traversal requirements

Performance metric

- Utilization of most congested link
- Bandwidth of business critical applications

Reformulating the problem

$$\max_{x \in X} \min_{y \in Y(x)} F(x, y)$$



LP dualization

$$\max_{\lambda, v, x} F'(\lambda, v, x)$$

Failure validation: Formulation

$$\max_{v, \lambda, x} \sum_{t, i \neq t} d_{it} (v_{it} - v_{tt})$$

$$s.t. \quad v_{it} - v_{jt} \leq \lambda_{ij} \quad \forall t, \langle i, j \rangle \in E$$

$$\sum_{\langle i, j \rangle \in E} \lambda_{ij} c_{ij} (1 - x_{ij}^f) = 1$$

$$x^f \in X$$

$$x_{ij}^f \in \{0, 1\}; \quad \lambda_{ij} \geq 0; \quad \langle i, j \rangle \in E$$

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Depends on failure model of interests

- E.g. simultaneous f link failures

$$\sum_{i, j} x_{ij}^f = f$$

Failure validation: Formulation

$$\begin{aligned} \max_{v, \lambda, x} \quad & \sum_{t, i \neq t} d_{it} (v_{it} - v_{tt}) \\ \text{s.t.} \quad & v_{it} - v_{jt} \leq \lambda_{ij} \quad \forall t, \langle i, j \rangle \in E \\ & \sum_{\langle i, j \rangle \in E} \lambda_{ij} c_{ij} (1 - x_{ij}^f) = 1 \\ & x^f \in X \\ & x_{ij}^f \in \{0, 1\}; \quad \lambda_{ij} \geq 0; \quad \langle i, j \rangle \in E \end{aligned}$$

Can be converted to mixed-integer linear program.
In general, validation problems could be non-linear.

Solution approach

- **Focus on upper bounds (relaxation)**
 - Intractable problems – hard to solve to optimality
 - Upper bounds sufficient for validation use
- **Goal: Develop a general approach**
 - Applicable to diverse validation problems (e.g., validating failures, demands...)
 - Yet, amenable to problem-specific structure
- **Use cutting-edge techniques from non-linear optimization**

Tractable relaxations: RLT

- RLT relaxations: general approach to relax non-convex problems into tractable LPs
 - Family of relaxations
 - Higher levels of hierarchy
 - Converge to optimal value of the non-convex problem
 - Incur higher complexity
- **For scalability, focus on the first level**

Our results on effectiveness of RLT

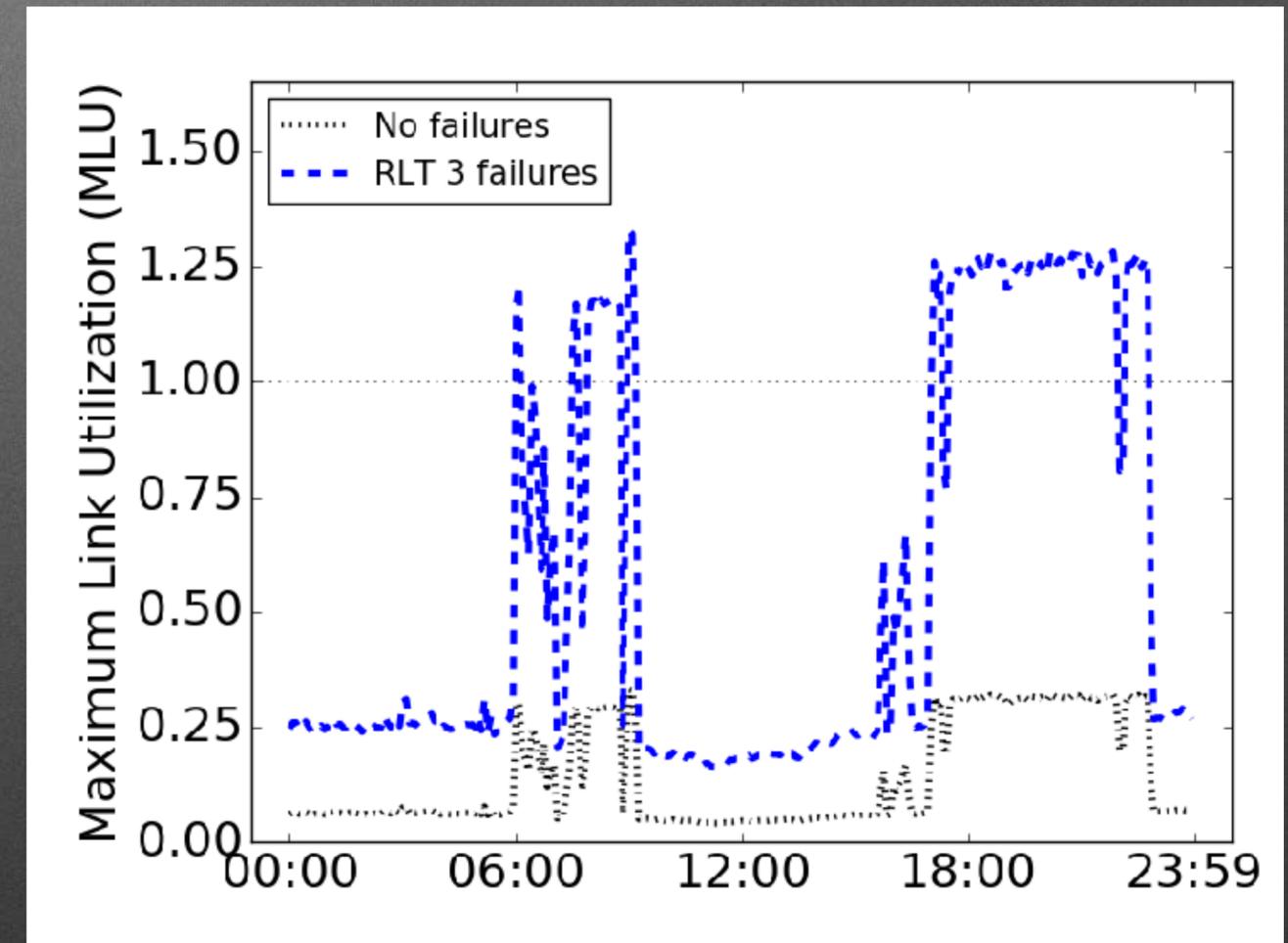
- Compare RLT with two theoretical benchmarks
 - Both bound worst case performance across failures/demands, but with limited network adaptation
 - Oblivious routing [Applegate, et al., Sigcomm '03; Wang, et al., Sigcomm '06, etc.]
 - Affine adaptation: a generalization of oblivious routing, studied in robust optimization
- **Our results show**
 - First-level RLT dominate oblivious/affine adaptations
 - Better results possible by exploiting problem-specific structure combined with RLT

Evaluation

- **Real topologies**
 - Abilene, GEANT, and ANS (from The Internet Topology Zoo)
- **Real and synthetic traffic matrices**
 - Real trace: 6-month end-to-end demand on Abilene
 - Synthetic: Gravity model

Results: Effectiveness of RLT

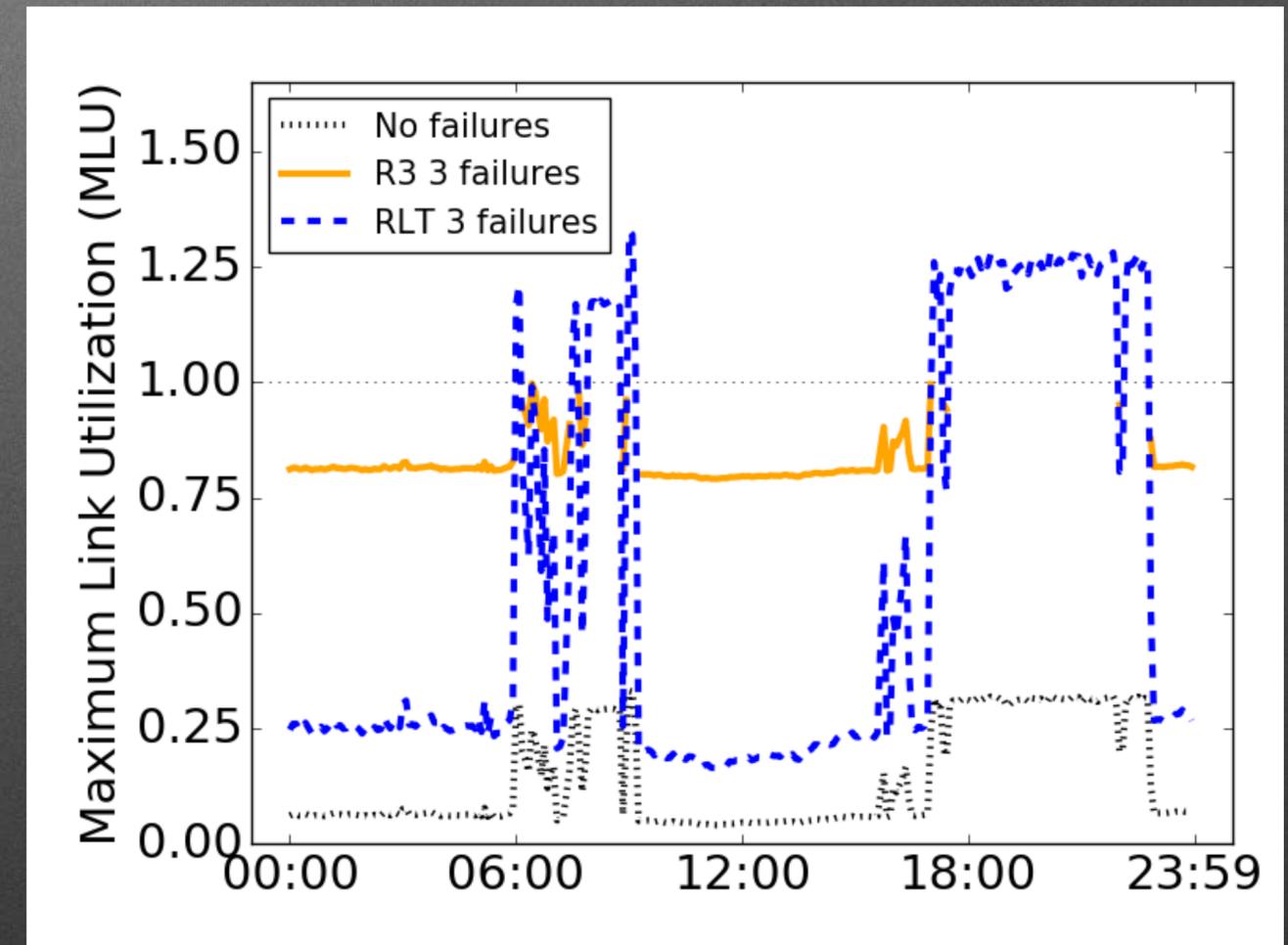
- Compare maximum link utilization (MLU)
- The optimal IP scheme vs. our RLT relaxation LP
- **RLT matches optimal** in all our experiments



Abilene Network — 3 link failures

Results: Effectiveness of RLT

- Compare with R3 [Wang et al., Sigcomm '10]
 - Determines if $MLU < 1$ under f failures
 - Gives a valid bound only when $MLU < 1$
 - Based on oblivious approach
- **Our result**
 - First-level RLT dominates R3 whenever R3 provides a valid bound
- Other advantages of our approach
 - Useful to detect bad failure scenarios, and the amount of exceeded link capacity
 - Generalizes to other validation problems



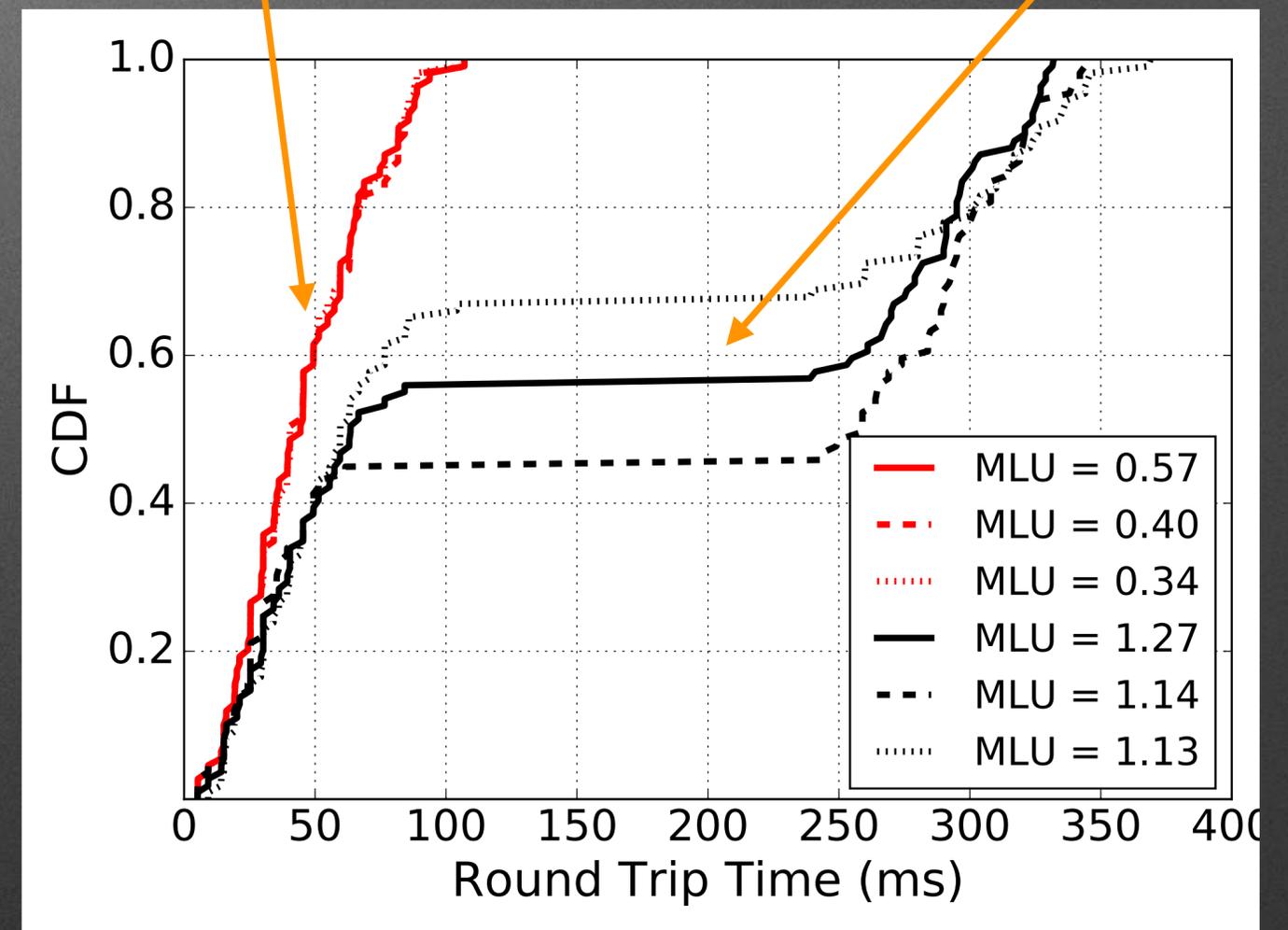
Abilene Network — 3 link failures

Using framework to detect bad failures

- Framework allows finding failures that impact the network the most
 - Random search not efficient
 - Only 0.05% of 3-failure scenarios are bad ($MLU > 1$)
- Emulate to understand latency behavior

Random scenarios

Bad scenarios



Emulated Abilene traffic matrix
with Mininet, and ONOS controller

Results: running time

- RLT relaxation LP vs. optimal IP (IP run for 2 hours)
- On scaled GEANT network (32 nodes, 1000 edges), 3 link failures:
 - RLT finished in **608** seconds, whereas IP finished in **3890** seconds
 - Only 60% of the IP instances completed in 2 hours
- Our RLT relaxation LP doesn't degrade with larger number of failures

Example: Tunnel selection validation

Uncertainty Set

- All f or fewer link failures
- Shared risk link group
- Weighted averages of historical demands

Adaptations

- Flexibly rerouting (Multi-commodity flow)
- Rerouting constrained to pre-selected tunnels
- Constrain with middlebox traversal requirements

Performance metric

- Utilization of most congested link
- Bandwidth of business critical applications

Problem:

- For a given choice of tunnels, are utilizations of all links across all traffic demands of interest within acceptable limits?

Tunnel selection: Results

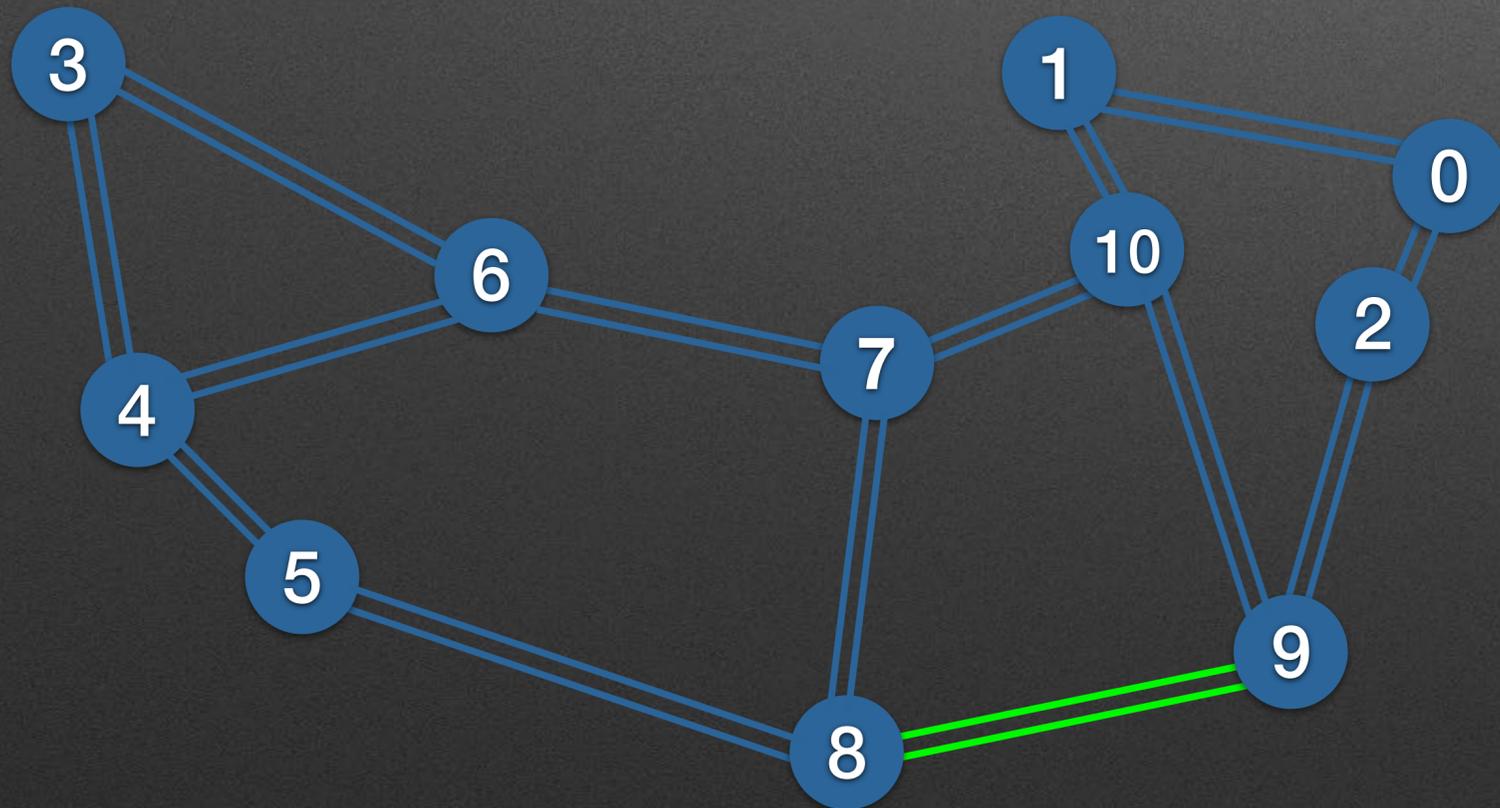
- **Predicted demand:** weighted averages of historical matrices
 - Validation problem is an LP
 - On Abilene: **First-level RLT achieves optimal MLU**
- **Widely-used tunnel selection heuristics may perform poorly**
 - E.g., K-shortest (SWAN, Sigcomm '13), Shortest-Disjoint heuristics
 - More robust tunnel selection heuristic performs much better

Synthesizing valid designs

- **Validation is a stepping stone for synthesis**
- **Example: Optimal Capacity Augmentation**
 - Incrementally add capacity to existing links
 - Minimizing cost of adding capacity
 - Ensure resulting network can handle all failure scenarios
- **One can use our framework for synthesis in 2 ways:**
 - 1) Get conservative solution, with a single LP
 - 2) Iterative approach, which gives a lower bound on cost at each step

Capacity augmentation: Abilene

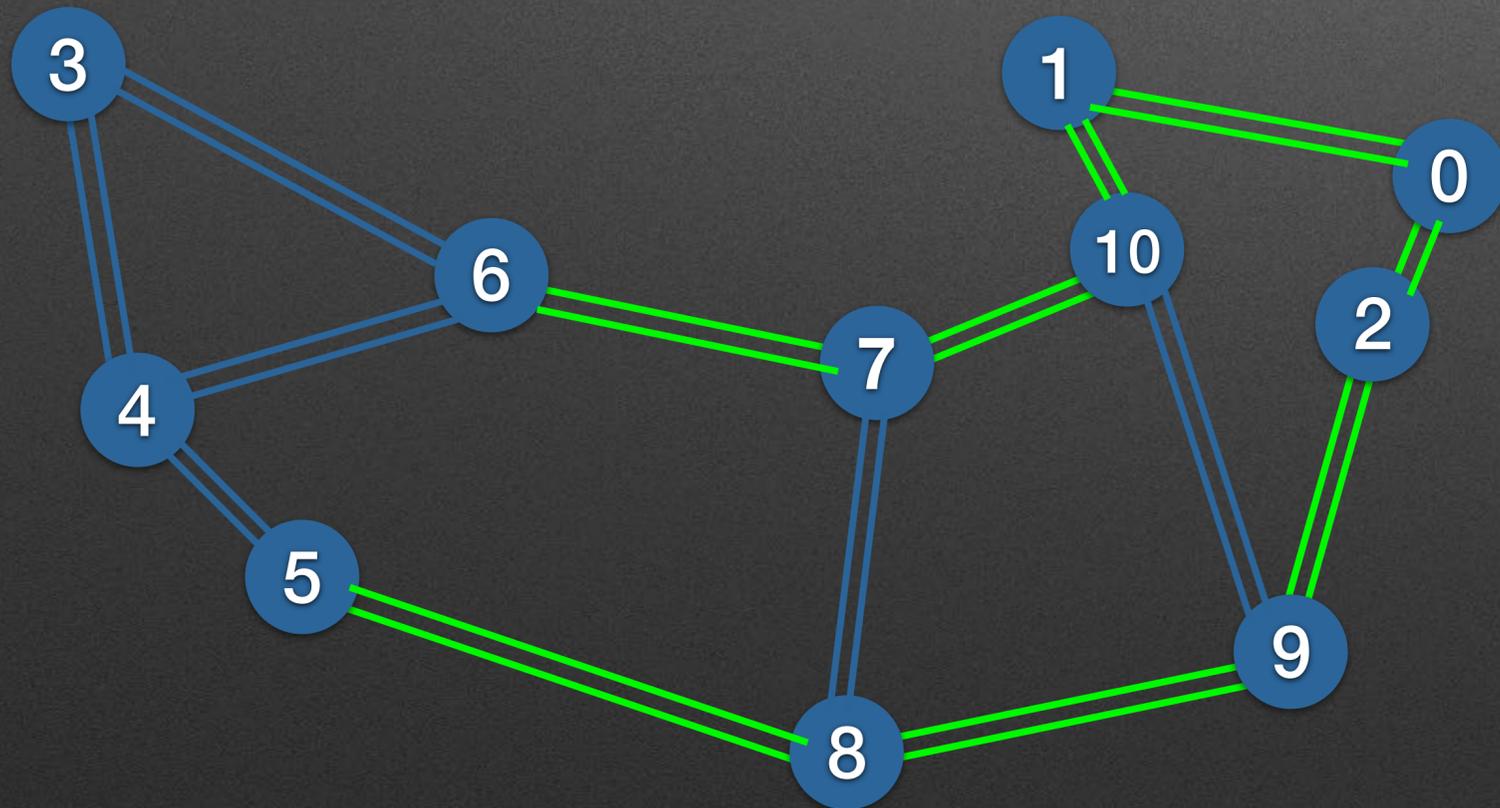
- Validate if $MLU \leq 1$.
- If not, run augmentation LP with counter examples



Step	Counter examples	MLU	Links Augmented	Total new capacity (Gbps)
1	(1, 10), (2, 9)	1.274	(1, 10)	2.744
2	(2, 9), (10, 1)	1.274	(2, 9)	5.488
3	(9, 8), (10, 7)	1.217	(9, 8)	7.653

Capacity augmentation: Abilene

- Validate if $MLU \leq 1$.
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2	(2, 9), (1, 10)	1.274	(2, 9)	5.488
3	(9, 8), (10, 7)	1.217	(9, 8)	7.653
4	(10, 7), (9, 8)	1.217	(10, 7)	9.818
5	(0, 2), (1, 10)	1.192	(0, 2)	11.743
6	(1, 0), (1, 10)	1.071	(1, 0)	12.452
7	(7, 6), (8, 5)	1.006	(7, 6)	12.509
8	(8, 5), (7, 6)	1.006	(8, 5)	12.566
9	—	1.000	—	—

Conclusions

- **Early effort at formally verifying quantitative network properties under uncertainty**
- **Generic framework for a wide class of network validation problems**
- **Modeling adaptivity results in intractable problems**
 - RLT relaxations promising
 - Tighter bounds than oblivious
 - Exact in multiple failures case and predicted demand case
- **Validation framework enables network synthesis**

Thanks!
Questions?